## Exploiting Abstract Interpretation for Model Checking Programs

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## Objective I

- Objective: to improve the performance of software model checking using static analysis techniques
- Static analysis and model checking: two formal verification techniques that can be used to verify that a program adheres to its specification
- Our new method of combining static analysis with model checking:
  - Has the potential to improve the performance of tools that use formal verification techniques
  - Is automatable, and we are currently working on its automation

## **Objective II**

- Basic outline of steps in our method:
  - Perform a static analysis of a C program
  - Insert the results of the analysis into the program
  - Run a model checking tool with the modified program as input

Test results: significant improvement in speed of the model checker using our method on some programs

## Our Approach

- Performing an abstract interpretation to identify variable values at varied program points
- 2 Using backward slicing to choose which program variables to track
- 3 Exploiting this information to reduce the search space of the model checker

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## Approaches to Program Verification

- Abstract Interpretation: An approximation of program semantics based on mappings between concrete and abstract lattices ⇒ symbolic evaluation in abstract domain
  - Usefulness of [nondeterministic, lossy] abstract program dependent on abstractions
  - Loops may require unrolling, with loss of precision (or an indeterminate fixed point computation)
- Model Checking with CEGAR: Iteration over abstraction model checking - refinement cycle to automatically prove program correctness
  - State space explosion
  - Success limited by choice of predicate abstractions

## **Existing Tools**

### Frama-C [2]:

- An extensible C verification framework
- Plugins include abstract interpretation (Eva) and slicing
- CPAchecker [1]:
  - Configurable program analysis dealing mainly with model checking of control-flow automata constructed from C programs
  - Includes support for CEGARish checking (in predicateAnalysis configuration)
- CegarMC [3]:
  - A previously published Frama-C plugin by the authors
  - Integrates CEGAR-based model checkers into Frama-C

- Two Frama-C plugins used by our method:
  - Eva to automatically compute sets of possible values for the variables of an analyzed program
  - **2** The program slicing plugin:
    - Reduces a program based on a backward slicing criterion, traditionally a program location and a set of program variables, so that the behavior of the original program is preserved with respect to the criterion
    - The results from Eva are also used to compute program slices

## Static Analysis Tools II

- Sets of possible values introduced into the program using assume statements
- Assume statements inserted at selected points throughout the program

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CPAchecker run with the modified program as input

## Example

```
int x, x0, y, y0;
y = 0;
while (x > 0) {
    x0 = x;
    y0 = y;
    x = x - 1;
    y = y + 2;
    if (2 * (x0 - x) != y - y0)
        error();
}
```

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```
int x, x0, y, y0;
v = 0;
assume(x >= INT_MIN && x <= INT_MAX &&
       xO >= INT_MIN && xO <= INT_MAX &&
       v == 0 \&\& v 0 == 0);
while (x > 0) {
  assume(x >= 1 && x <= INT_MAX &&
         xO >= INT_MIN && xO <= INT_MAX &&
         y >= 0 \&\& y <= INT_MAX - 1 \&\&
         y % 2 == 0 &&
         yO >= 0 \&\& yO <= INT_MAX - 1 \&\&
         v0 \% 2 == 0);
  x0 = x;
  v0 = v;
```

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## TOOL DEMO

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## Frama-C Architecture I



From Frama-C Plugin Manual

### Plugins:

- Interfaces to abstract syntax tree (AST), C intermediate language (CIL), AI lattices, etc. provided by kernel
- Plugins used for either analysis (≥ 1 AST) or source-to-source transformation (> 1 AST)
- Kernel-integrated plugins include Eva and wp (statically linked)

## Frama-C Architecture II



From Frama-C Plugin Manual

#### Plugins:

- Extensible through user-written plugins, typically linked dynamically
- Common plugin interface allows for inter-plugin information sharing, along with a central mechanism for combining results

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All programmed in OCAML

## **Tool Architecture**



From Frama-C Plugin Manual

- Fine-tune the abstract interpretation
  - E.g., the program points where the abstract interpretation information is exploited
- Evaluate our method with more and a wider variety of example programs
- Fully automate our method, extending our CegarMC plugin

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- Explore other possible combinations of abstract interpretation and model checking:
  - **Residual program**: unexplored parts of a model check
  - Create a residual program generator using Frama-C plugins
  - Pass the residual program generated by CPAchecker along with any other necessary information to a Frama-C plugin

## References

- [1] D. Beyer and M. E. Keremoglu.
  - CPAchecker: A tool for configurable software verification. In *Computer Aided Verification (CAV)*, pages 184–190, 2011.
- [2] F. Kirchner, N. Kosmatov, V. Prevosto, J. Signoles, and B. Yakobowski.

Frama-C, a software analysis perspective.

In *Formal Aspects of Computing*, 27: pages 573–609, March 2015.

[3] S. Shankar and G. Pajela.

A Tool Integrating Model Checking into a C Verification Toolset.

In International Symposium on Model Checking Software, pages 214–224, 2016.

# Thank you!

# Questions?

Frama-C: downloadable from www.frama-c.com CPAchecker: downloadable from cpachecker.sosy-lab.org CegarMC Plugin: downloadable from http://www.compsci.hunter.cuny.edu/~sshankar/cegarmc.html